

Logic Sniffer Specification (Demon Core FPGA)

Version 1.03

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Revision History

Feb 23, 2011 - Rev 1.01 - Initial Release.

May 7, 2011 - Rev 1.02 - Clarified various points/problems (much thanks to flubberlab).

Jan 1, 2012 - Rev 1.03 - Added section on implementing inverted trigger inputs.

1.1 Advanced Trigger

The advanced trigger offers many features found in an HP16500 / 16550 timing logic analyzer. It offers 16 sequence levels with two way branching (a programmable state machine), ten pattern value comparisons, two range checks, two edge checks, and two counter/timers.

There are several different modules within the advanced trigger, each needing its own spin on configuration. These are:

- **Trigger Terms.** Terms perform the actual bitwise masked pattern comparisons on the 32-bit incoming data. There are ten terms. Each one needs 128 bits of configuration data.
- **Trigger Range Detectors.** The two range detectors look for values falling between a lower & upper 32-bit limit. Each range check uses two 32-bit magnitude comparisons., one for the upper compare, one for the lower. A match means: “upper \geq indata \geq lower”. Each range detector needs 512 bits of configuration data.
- **Trigger Edge Detectors.** The two edge detectors compare sampled input against a delayed version. Rising edges, falling edges, or both, or neither can thus be detected. Each needs 256 bits of configuration data.
- **Trigger Timers.** Timers are started/stopped under control of the trigger sequence states. There are two timers, each needing a 36 bit limit value. The timers have a range from 10ns to over 11 minutes.
- **Trigger Sums.** These combine the results of the trigger terms, range comparisons, edge checks, and timer checks. There are 11 logic units programmable to perform any logic function. Examples are AND/NAND/OR/NOR/XOR/NXOR. The logic units are combined in a tree structure, and combine/sum all trigger sources down to a match/miss signal. There are three trigger sums per trigger state, and 16 trigger states, totaling 48 sums. Each sum needs 192 bits of configuration data.
- **Trigger Sequence States.** Each state is guided by the results of three trigger sums (see above). If the “hit” trigger sum matches, for a specified number of occurrences, the FSM advances. A hit also controls if a timer is started/stopped/cleared. If the “else” trigger sum matches, the FSM branches to the “else” state. Data capture is controlled by the “capture” trigger sum – if it matches, a sample is taken. If not, then not. Each state takes 32 bits of configuration.

Each module will be described in detail in the following chapters.

2 Background

A little background to help put things in context...

2.1 FPGA's

The FPGA (field programmable gate array) is very useful for creating custom hardware. They can implement anything from simple logic up to entire computers. To accomplish this, an fpga contains large numbers of tiny lookup tables and flip-flops as building blocks, plus usually some larger memories (SRAM's).

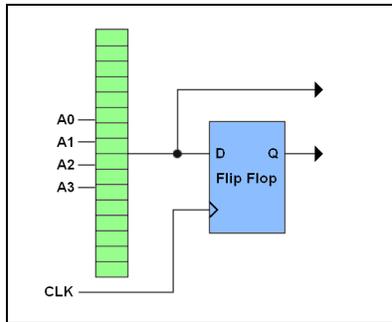


Figure 2 - Logic Block

Add some I/O pins for talking to the outside world, plus lots of configurable wires to connect everything, and you have an FPGA.

The lookup tables are the interesting piece, and store only 16 bits of data. They have four inputs (the address), an output, and some config pins. Normally they are configured to form logic. Any four-input function can squeeze in there.

For example, a four input AND gate. Load LUT addresses 0 to 14 with all zeros, and set address 15 to one. Now if all inputs are set, you get a one on output -- otherwise zero. Just like an AND gate.

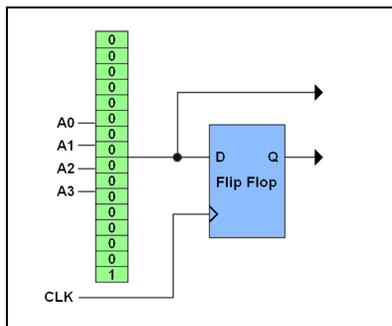


Figure 4 - AND gate equivalent

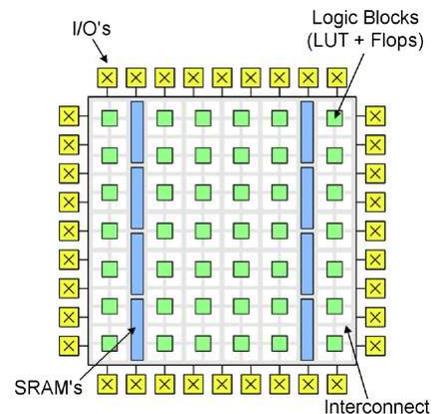


Figure 3 - FPGA Layout

In the case of a logic analyzer, even more useful functions are possible. For example, a trigger search on four bits of sampled data. The LUT allows easy matching of any pattern.

LUT's are normally programmed serially when the FPGA powers up. Think shift register. However, they can also be selectively changed during runtime.

2.2 The Datapath

The Logic Sniffer is composed of several main blocks. Most are concerned with the data path (see right). In the drawing thin lines are control signals. Thick lines are the data path.

The “*sync module*” captures input data using the sample clock. The sample clock can either be the internal 100Mhz reference clock, or supplied from an external source.

The “*async fifo*” transfers the captured input to the rest of the fpga. If the sample clock & internal reference clock are totally unrelated, the fifo avoids timing problems.

The “*sampler module*” is what implements the “sampling rate”. Assume the sample clock is using the 100Mhz reference. If you want a 10Mhz sample rate, the sampler forwards every tenth valid capture. For 1Mhz, every hundredth valid capture, etc...

The “*triggers*” look for patterns. They control what ultimately gets stored in SRAM -- with the “*capture*” signals -- and tell the controller when a trigger hit occurs. The basic trigger captures everything always. The advanced trigger lets you be more selective.

The “*delay fifo*” aligns its output to the capture/run trigger outputs. The trigger modules take a few clocks to evaluate things. For a trigger “*capture*” to work, the input to the next state must match.

The “*data align*” module removes gaps within the input data caused by disabled groups. If you decide to disable groups 0 & 2, the aligner shifts the remaining groups down to fill the gap. This simplifies things later.

The “*RLE encoder*” looks for repeat captures. If found, counts are stored instead of repeats of unchanging data. This can greatly increase the effective storage capacity.

The “*controller*” is what decides how much gets captured, and when to send results to the client on your PC.

The “*SRAM interface*” handles saving & fetching captured data to & from the SRAM memory.

The “*SPI interface*” uploads the captures serially to the client. Only valid data is sent. Gaps due to disabled groups are filtered out.

A few remaining blocks aren't shown here. The inbound SPI interface for receiving commands, the command decoder & flags register.

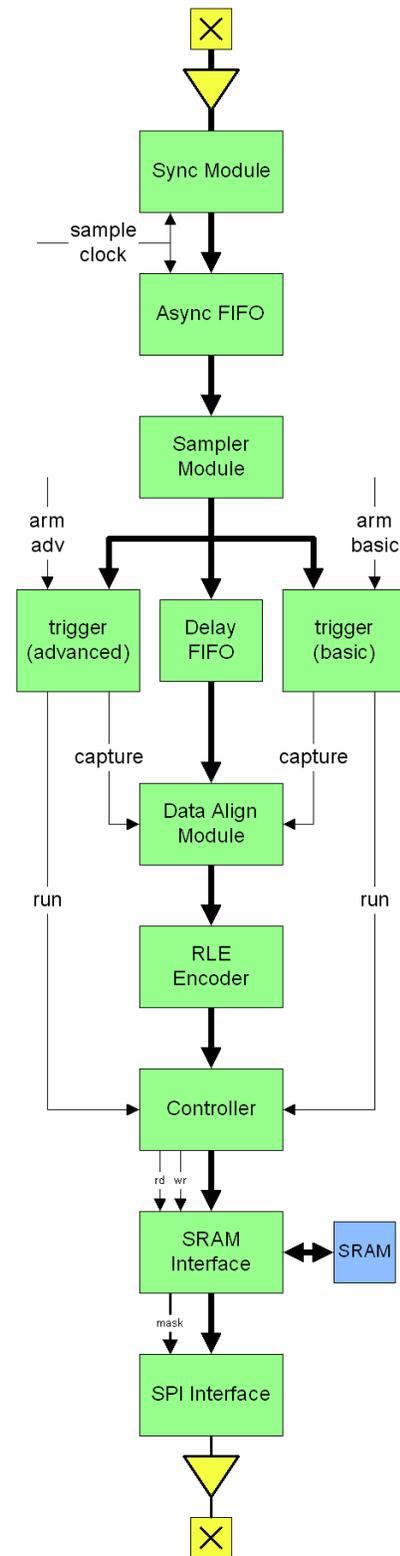


Figure 5 - Datapath Diagram

2.3 The SUMP Protocol

The SUMP communication protocol assumes a standard RS232 port connection. Although the Logic Sniffer is USB based it emulates a very fast serial port from the perspective of the PC.

When sending captured data the analyzer sends back a contiguous block of data. Disabled channels are not sent however. First byte sent corresponds to sampled bits 7:0 (assuming channel 0 isn't disabled). The most recently captured data is returned first (ie: data is sent to client in reverse order). If Demux mode is enabled, group 2 and 3 inputs are ignored. Channel group 0 and 1 inputs are instead captured (if enabled) twice per sample - once each clock edge - and returned as a pair of bytes per group.

The following list documents the commands understood by the Logic Sniffer FPGA.

2.3.1 Short Commands (one byte commands)

0x00 - Reset

Reset the analyzer. Send at least five times to ensure FPGA in a known state. Clears any previous long commands that might have gotten stuck. The reset command should be issued prior to each Arm command (basic/legacy or advanced), to ensure previous captures/triggers are cleared.

0x01 - Arm Basic/Legacy Trigger

Arm the basic trigger. Sampling begins immediately. Once the trigger fires, the controller waits for "delay count" additional samples before returning captured data to client.

0x02 - Query ID

The device responds with four bytes. Currently "SLA1" to maintain backwards compatible SUMP protocol, output LSB first. ie: "1", "A", "L", "S"

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0x04 - Query Meta Data

The device responds with a variable length block of data composed of one or more tokens + associated data.

- Token 0 denotes end of meta data
- Tokens 0x01 to 0x1F are followed by a null terminated string (UTF-8 encoded).
- Tokens 0x20 to 0x3F are followed by a 32-bit unsigned integer (MSB first).
- Tokens 0x40 to 0x5F are followed by a 8-bit unsigned integer.
- All other token ranges are reserved.

The following fields are currently defined:

| Token | Description | FPGA Currently Returns |
|-------|----------------------------------|----------------------------|
| 0x00 | End of Meta Data | 0 |
| 0x01 | Device Name | "Open Logic Sniffer v1.01" |
| 0x02 | Version of FPGA firmware | "3.0" |
| 0x03 | Ancillary Version (PIC firmware) | - |
| 0x20 | Number of Probes | - |
| 0x21 | Sample memory available (bytes) | 24576 |
| 0x22 | Dynamic memory available (bytes) | - |
| 0x23 | Maximum sample rate (Hz) | 200000000 |
| 0x24 | Protocol version | - |
| 0x25 | Capability Flags | 0x0000001F |
| 0x40 | Number of Probes (short version) | 32 |
| 0x41 | Protocol version (short version) | 2 |

Clients should accept data from either long or short versions of a meta field, if both are possible.

The Capability Flags provide a bitmask indicating available FPGA features. Moving forward, this should help decouple the fpga from a needing a specific version of client (or visa versa).

- Bit 0 - Basic Trigger available.
- Bit 1 - Advanced Trigger available.
- Bit 2 - RLE Encoding available (see flag register).
- Bit 3 - Extra RLE Encoding modes available (see flag register).
- Bit 4 - State capture mode available (see flag register)
- Bit 5 - Finish Now command available.
- Bit 6 - Query Input Data command available.
- Bit 7 - Query Capture State command available.
- Bit 8 - Return Capture Data command, and manual capture mode available (see flag register).
- Bits 31-9 - reserved

0x05 - Finish Now (disable RLE mode)

Turns off RLE mode specified in the flags register (see 0x82 command below). In RLE mode, the analyzer can sit forever if inputs aren't changing. This gives the PC client application an opportunity to finish a capture on demand.

Finish now also aborts the advanced trigger, & forces data capture to quickly fill the buffer & end capture.

0x06 - Query Input Data

The device responds with four bytes. Gives a snapshot of the current logic analyzer input bits.

0x07 - Poll/Query Analyzer State

Returns capture state data. Uses same encoding as meta data, with tokens & parameters. ie:

- Token 0 denotes end of meta data
- Tokens 0x01 to 0x1F are followed by a null terminated string (UTF-8 encoded).
- Tokens 0x20 to 0x3F are followed by a 32-bit unsigned integer (MSB first).
- Tokens 0x40 to 0x5F are followed by a 8-bit unsigned integer.
- All other token ranges are reserved.

| Token | Description |
|-------|----------------------------------|
| 0x00 | End of Poll Data |
| 0x30 | Flags[31:0] |
| 0x31 | Capture Depth |
| 0x32 | Trigger Position |
| 0x50 | Basic Trigger Level (0 to 3) |
| 0x51 | Advanced Trigger Level (0 to 15) |

Flags:

- bit 0 - Analyzer armed.
- bit 1 - Capturing pretrigger data (used in timing mode only).
- bit 2 - Capture complete.
- bit 3 - Capture cancelled.
- bit 4 - Basic trigger active. Basic trigger level (0x50 field) valid.
- bit 5 - Basic trigger has fired.
- bit 6 - Advanced trigger active. Advanced trigger level (0x51 field) valid.
- bit 7 - Advance trigger has fired.
- bits 31-8 - reserved

0x08 - Return Capture Data

If the fpga is in "manual capture mode" (see flag register), this command returns any capture data to the client. This command also cancels any active capture. Note: Only data captured since the fpga was armed is returned. If there were only 8 captures, the client only gets 8 captures. Client should query capture state to obtain trigger position.

0x0F - Arm Advanced Trigger

Arm the advanced trigger. Conditional capturing (under control of trigger sequencer) begins immediately. Once trigger fires, the controller waits for "delay count" additional samples before returning captured data to client.

0x11 - XON

Deprecated. On a serial port, this would pause output when FPGA was returning captured data. Currently ignored by the FPGA.

0x13 - XOFF

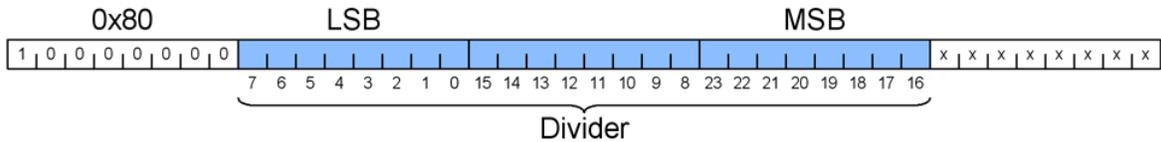
Deprecated. On a serial port, this would resume output when FPGA was returning captured data. Currently ignored by the FPGA.

2.3.2 Long Commands (five byte commands)

0x80 - Set Sample Rate

The Logic Sniffer samples at rates lower than the reference clock by using a *divisor-count*. Every given number of clocks, a sample is accepted. To set the sample rate, the divisor count is computed as follows:

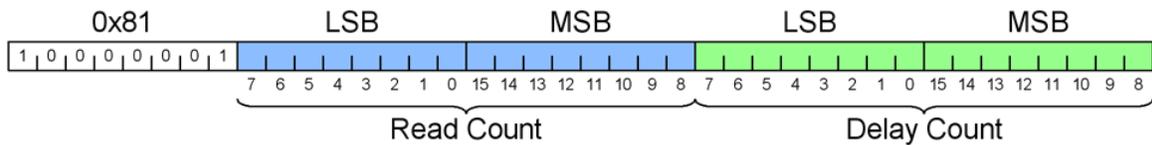
$$\text{divider} = (\text{clock} / \text{sample-rate}) - 1$$



0x81 - Set Read Count & Delay Count

The *read-count* specifies the number of samples (divided by four) to be returned to the client on your PC. The *delay-count* is the number of samples (divided by four) to capture -after- the basic or advanced trigger fires. Samples are defined as one byte per each enabled input group.

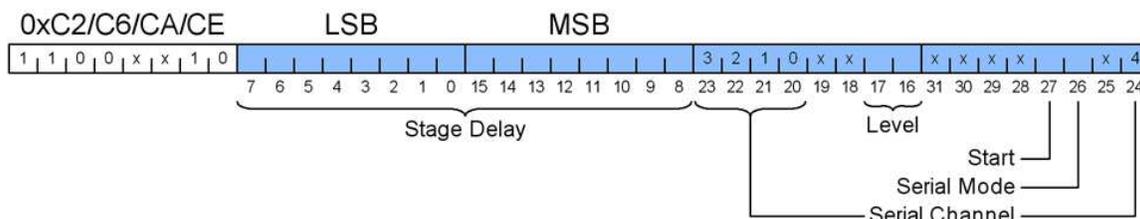
A *read-count* larger than the *delay-count* returns captures from before the trigger. Such pre-trigger captures are only valid if the device was running long enough before the trigger fired. ie: a trigger that fires immediately could return garbage.



0xC2/C6/CA/CE - Set Basic Trigger Configuration

Writing this register configures the selected basic trigger stage. There are various fields available for customizing the operation of the trigger:

- **Stage Delay.** If a trigger match occurs, the action of the stage is delayed by number of samples.
- **Serial Mode.** If enabled, trigger operates as a serial trigger. 1 = Serial Mode.
- **Serial Channel.** In serial mode, select which input pin to treat as serial channel (0 to 31).
- **Level.** Trigger level at which current stage will become active. Level 0 means activate immediately.
- **Start.** Tell controller the basic trigger has fired.



Serial Mode:

In serial mode, a basic triggers stage monitors only a single sampled input bit, selected with the "serial-channel". The selected bit is fed into a serial-to-parallel shift register. Thus the last 32 samples of a single bit can be evaluated using the normal masked compare logic.

To be effective, you should use an external clock in serial mode. Otherwise it's very difficult to predict exactly when samples are taken & added to the shift register. So for example, when monitoring a SPI bus use the SPI clock for capturing.

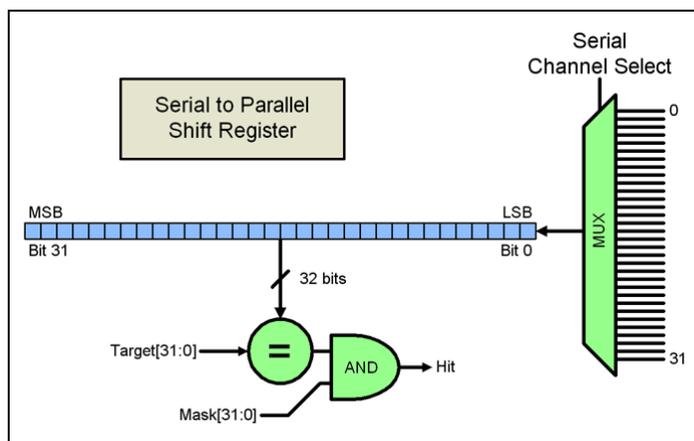


Figure 6 - Basic Trigger Serial Mode

Trigger Levels:

The basic trigger has a simple level counter. When any trigger stage matches something, the level count is incremented. When the level count matches a given stages "trigger-level", the stage becomes active. Thus you can search for a simple sequence of events.

ie: Setup trigger 0 to start immediately, triggers 1 & 2 to start on level 1, and trigger 3 to start on level 2. When trigger 0 matches, both triggers 1 & 2 activate. When either of them matches, trigger 3 activates.

3 Trigger Terms

A trigger term is pretty simple. It compares the logic analyzer input against a target. If the input matches the target, it signals a hit. Trigger terms work the same in both basic & advanced triggers.

To make things a little more useful, the fpga also supports masking (or ignoring) parts of the input. The mask determines which bits must match, so you only compare them & nothing else.

As shown below, the compare is done with an XNOR (exclusive inverted or), which goes high when both inputs match. Masking is done by an AND gate.

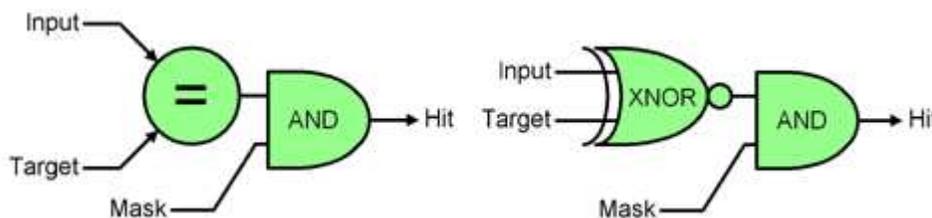


Figure 7 - Masked Compare

To implement something like this normally requires flops to hold the target, flops to hold the mask, and logic for the XNOR & AND gates. So around 64 flops plus 32 logic blocks.

However, a better way exists. In an fpga, LUT's can handle everything. The 32-bit target & mask values, and logic are all encoded directly into only eight LUT's.

Quite a savings, but there is a downside. Programming the LUT is non-trivial.

The "complexity" of the masked compare doesn't go away -- its simply transferred elsewhere. Thus in the advanced trigger, the eight LUT's need 128 bits of configuration supplied by the client.

The LUT memory is programmed with the -results- of a masked compare. The target & mask are known, and the input sample is one of sixteen possible values.

Thus we simply run through all combinations in advance. ie:

```
LUT address 0 stores result of: (0 ^ value) & mask == 0
LUT address 1 stores result of: (1 ^ value) & mask == 0
LUT address 2 stores result of: (2 ^ value) & mask == 0
LUT address 3 stores result of: (3 ^ value) & mask == 0
LUT address 4 stores result of: (4 ^ value) & mask == 0
etc...
```

| Mask | Input | Target | Hit |
|------|-------|--------|-----|
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

Figure 8 - Truth Table

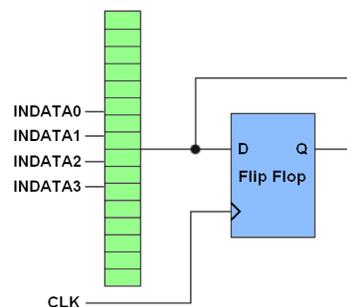
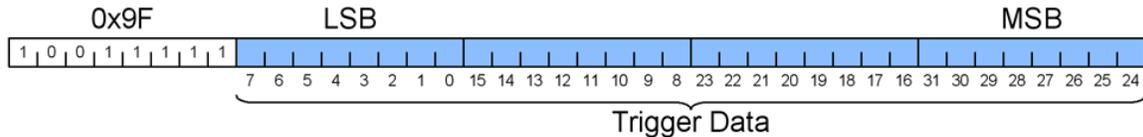


Figure 9 - Trigger Term LUT

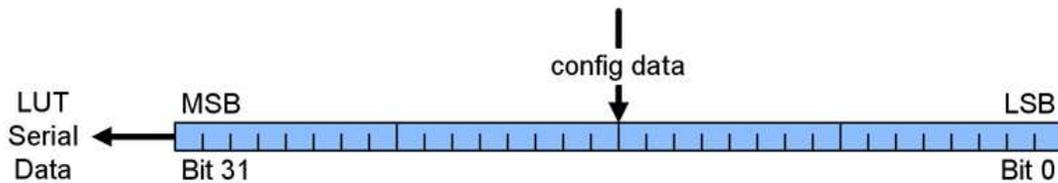
For each of the eight LUT's, the above is repeated with different 4 bit chunks of value & mask.

3.1 Programming Lookup Tables

To enable easier LUT programming, the fpga contains a built in parallel-to-serial converter. You simply write a long SUMP command to the chain address. Note the SUMP protocol dictates the least significant data byte be output first.



Once the fpga receives the long command, the bytes are organized MSB first (bits 31 down to 0), and then serially shifted into the LUT chain.



Some of the advanced trigger modules are quite complex and need multiple writes to fully load their LUT chains. Chains vary in length up to 512 bits.

3.2 Example (Trigger Term Initialization)

```
#define write_select(value) write_long_command (0x9E, value)
#define write_chain(value) write_long_command (0x9F, value)

void write_term (int term_number, unsigned int target, unsigned int mask)
{
    unsigned int bitmask=1;
    unsigned int lutvalue0=0;
    unsigned int lutvalue1=0;
    unsigned int lutvalue2=0;
    unsigned int lutvalue3=0;

    for (int i=0; i<16; i=i+1) {
        if (((i ^ (target & 0xF)) & (mask & 0xF))==0) lutvalue0 |= bitmask;
        if (((i ^ ((target>>4) & 0xF)) & ((mask>>4) & 0xF))==0) lutvalue0 |= (bitmask<<16);
        if (((i ^ ((target>>8) & 0xF)) & ((mask>>8) & 0xF))==0) lutvalue1 |= bitmask;
        if (((i ^ ((target>>12) & 0xF)) & ((mask>>12) & 0xF))==0) lutvalue1 |= (bitmask<<16);
        if (((i ^ ((target>>16) & 0xF)) & ((mask>>16) & 0xF))==0) lutvalue2 |= bitmask;
        if (((i ^ ((target>>20) & 0xF)) & ((mask>>20) & 0xF))==0) lutvalue2 |= (bitmask<<16);
        if (((i ^ ((target>>24) & 0xF)) & ((mask>>24) & 0xF))==0) lutvalue3 |= bitmask;
        if (((i ^ ((target>>28) & 0xF)) & ((mask>>28) & 0xF))==0) lutvalue3 |= (bitmask<<16);
        bitmask <<= 1;
    }

    // write data into LUT serial chain. MSB must go in first. Total of 128 bits.
    write_select (0x20 + (term_number%10));
    write_chain (lutvalue3);
    write_chain (lutvalue2);
    write_chain (lutvalue1);
    write_chain (lutvalue0);
}
```

Figure 10 - Initialize Trigger Term LUT Values

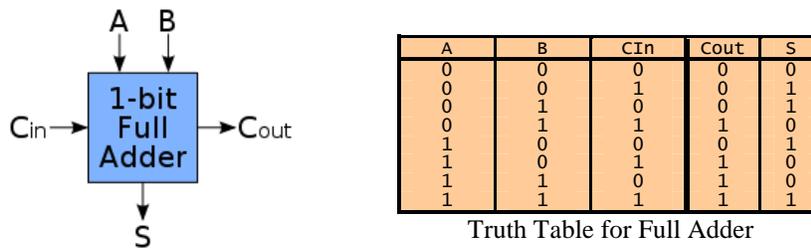
4 Range Detectors

The advanced trigger supports detecting when sampled data falls within a range of values. ie: if the sample is between a lower limit & upper limit. ie:

$$\text{hit} = (\text{upper-limit} \geq \text{sample}) \ \&\& \ (\text{sample} \geq \text{lower-limit})$$

So we need magnitude comparators. In hardware, that means adding or subtracting two numbers to see if something overflows or "carry's" over. A bit like long addition. Add a column, carry the one, add another, etc...

In hardware there is parallel to long addition: the full adder. The adder adds $A+B+C_{in}$, and outputs S & C_{out} , where all signals are single bits. To add wide values you chain the carry signals as needed.



Chains of full adders work, but are slow. The carry-out signal is derived from the inputs. In a 32-bit adder, there would be 32 chained carry circuits. Achieving timing would be harder than necessary.

An improvement is the carry-look-ahead adder. There is still a combinatorial element per bit, but it's very fast -- a simple mux. The fpga even has an integral "fast-carry-chain" circuit available for implementing such adders.

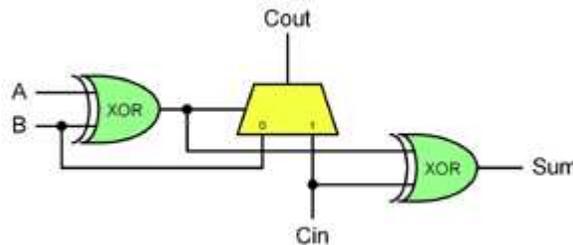


Figure 11 - Carry Look Ahead Adder

To make things even better, a magnitude comparison doesn't need the sum output at all. Just the fast carry chain. Thus we end up with something that fits perfectly into a single Logic Block. The LUT table simply needs loading with the result of an XOR between the input on A0, and the target value.

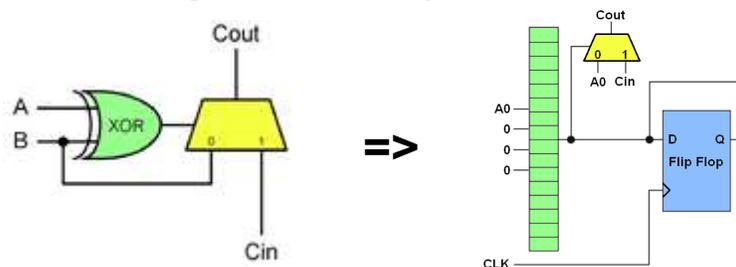


Figure 12 - Carry Look Ahead Comparator

4.1 Complex Range Compares

The fast carry change logic has a couple of interesting possibilities.

First of all, the range comparison can start anywhere. If the first bit of interest is bit 7 (for example), simply fill all lower LUT's with 1's. This NOP's them in the comparator.

Second, the bits participating in the range comparison don't need to be contiguous. You could setup to range check the value formed by indata bits 7,9,11, & 15. Simply NOP the LUT's of the don't care indata bits.

4.2 Prepare Target Values for Comparator

There are two ways to compare a value. Subtract a value & look for underflow/borrow, or add & look for overflow/carry. Since we have the fast-carry-chain, we're using the latter approach.

A given range target must be bitwise inverted before being programmed. For non-contiguous range checks, the inverted value must also be spaced out, across only those bits of interest.

Lower value = ~(lower_target - 1)

Upper value = ~(upper_target)

Yields a hit if (upper_target >= indata >= lower_target)

4.3 Range Compare LUT Chains

There are 32 LUT's in each range-check magnitude comparator. Each needs 16-bits, so that totals 512 bits of configuration. Two LUT's are programmed at once for each write to the chain-data register.

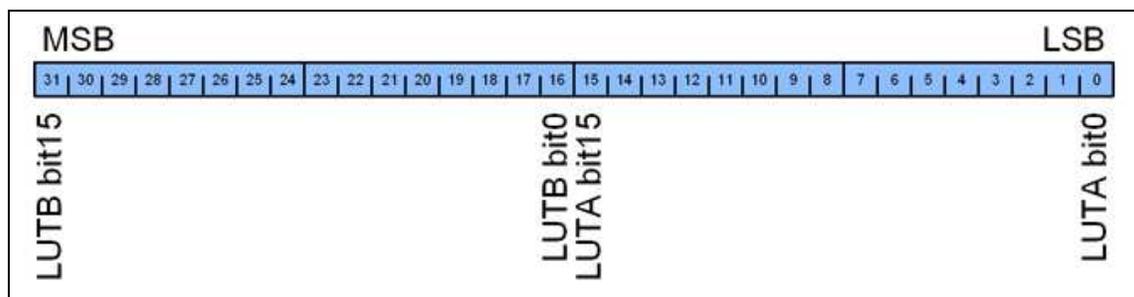


Figure 13 - Range LUT Organization

4.4 Example (Range Detect Initialization)

```
#define RANGE_XOR0 0xAAAA
#define RANGE_XOR1 0x5555
#define RANGE_NOP 0xFFFF

//
// Setup LUT's for Range Detectors
//
// Inputs:
// rangesel: 0=range1-lower, 1=range1-upper, 2=range2-lower, 3=range2-upper
// target:   Limit value.
// mask:     Indicate which bits should participate in range compare.
//
void write_range (int rangesel, unsigned int target, unsigned int mask)
{
    unsigned int value;
    unsigned int lutvalue=0;
    int i;

    write_select(0x30+(rangesel&3));

    // Count # of bits in mask...
    int bitcount=0;
    unsigned int bitmask=1;
    for (i=0; i<32; i++, bitmask<<=1)
        if (mask & bitmask)
            bitcount++;

    // Prepare target value...
    if (rangesel & 1)
        value = ~target; // upper range target
    else value = ~(target-1); // lower range value

    // Push MSB of target into bit 31...
    value <<= (32-bitcount);

    // Generate & program LUT values. Total of 512 bits.
    for (i=0; i<16; i=i+1) {
        if (((mask>>31)&1)==0)
            lutvalue = RANGE_NOP;
        else {
            lutvalue = ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
            value <<= 1;
        }
        mask <<= 1;
        lutvalue <<= 16;

        if (((mask>>31)&1)==0)
            lutvalue |= RANGE_NOP;
        else {
            lutvalue |= ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
            value <<= 1;
        }
        mask <<= 1;
        write_chain(lutvalue);
    }
}
```

Figure 14 - Initialize Range Detector LUT Values

5 Edge Detectors

The edge detectors compare a delayed sample to the current one. If any bit changes, with a rising edge, falling edge, either rising or falling edge, or neither (no change), it can be detected.

Each 4-input logic block evaluates two bits of input, and two bits of delayed input.

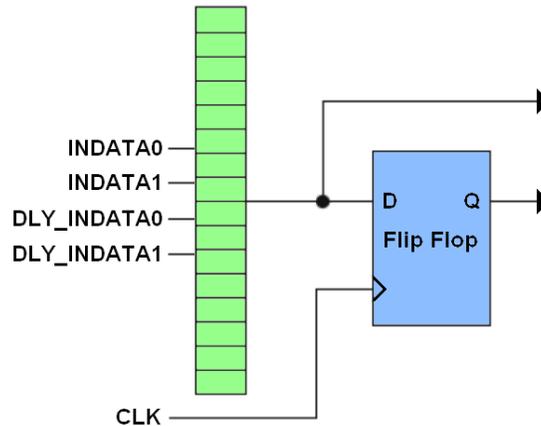


Figure 15 - Edge Detector LUT

All that is left now is placing 1's in the LUT where the edges of interest occur. Since two bits are evaluated in parallel, this becomes slightly more complex.

The following truth table shows the possible combinations.

| DLY INDATA1 | DLY INDATA0 | INDATA1 | INDATA0 | RISE0 | FALL0 | RISE1 | FALL1 | BOTH0 | BOTH1 | NONE0 | NONE1 |
|----------------|----------------|---------|---------|-------|-------|-------|-------|-------|-------|-------|-------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |

Rising edges are detected when DLY_INDATA is zero, and INDATA is one. Falling edges when DLY_INDATA is one, and INDATA is zero. Both edges are detected by simply combining columns. Neither edge is detected by inverting the "both" column.

5.1 Example (Edge Detector Initialization)

```
#define EDGE_RISE0 0x0A0A
#define EDGE_RISE1 0x00CC
#define EDGE_FALL0 0x5050
#define EDGE_FALL1 0x3300
#define EDGE_BOTH0 (EDGE_RISE0|EDGE_FALL0)
#define EDGE_BOTH1 (EDGE_RISE1|EDGE_FALL1)
#define EDGE_NEITHER0 (~EDGE_BOTH0 & 0xFFFF)
#define EDGE_NEITHER1 (~EDGE_BOTH1 & 0xFFFF)

void write_edge (int edgesel,
                unsigned int rising_edge, unsigned int falling_edge, unsigned int neither_edge)
{
    write_select (0x34 + (edgesel&1));

    unsigned int lutvalue=0;
    unsigned int bitmask = 0x80000000;
    for (int i=0; i<16; i=i+1) {
        // Evaluate indata bit1..
        if (neither_edge && bitmask)
            lutvalue |= EDGE_NEITHER1;
        else {
            if (rising_edge & bitmask) lutvalue |= EDGE_RISE1;
            if (falling_edge & bitmask) lutvalue |= EDGE_FALL1;
        }
        bitmask >>= 1;

        // Evaluate indata bit0..
        if (neither_edge && bitmask)
            lutvalue |= EDGE_NEITHER0;
        else {
            if (rising_edge & bitmask) lutvalue |= EDGE_RISE0;
            if (falling_edge & bitmask) lutvalue |= EDGE_FALL0;
        }
        bitmask >>= 1;

        if ((i&1)==0)
            lutvalue <<= 16;
        else {
            write_chain (lutvalue); // write total of 256 bits
            lutvalue = 0;
        }
    }
}
```

Figure 16 - Initialize Edge Detector LUT Values

6 Timers

Timer limits are very simple. You write a 36 bit value. Lower 32 bits to trigger register 0x38/0x3A. Upper 4 bits to trigger register 0x39/0x3B. The fpga reference clock is 100Mhz, so each timer tick corresponds to 10ns.

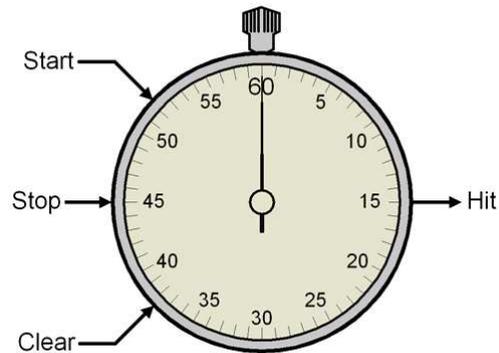


Figure 17 - Timer

The 36-bit timers have a range from 10ns to 687 seconds (over 11 minutes). Timers are started, stopped, and cleared under control of the trigger sequencer.

6.1 Example (Timer Initialization)

```
void write_trigger_limit (int timersel, uint64_t value)
{
  write_select (0x38 + (timersel&1)*2);
  write_chain (value & 0xFFFFFFFF);
  write_select (0x39 + (timersel&1)*2);
  write_chain (value>>32);
}
```

Figure 18 - Initialize Timer Values

7 Trigger Sums

The trigger summing logic combines the results of the trigger terms, range detectors, edge detectors and timers.

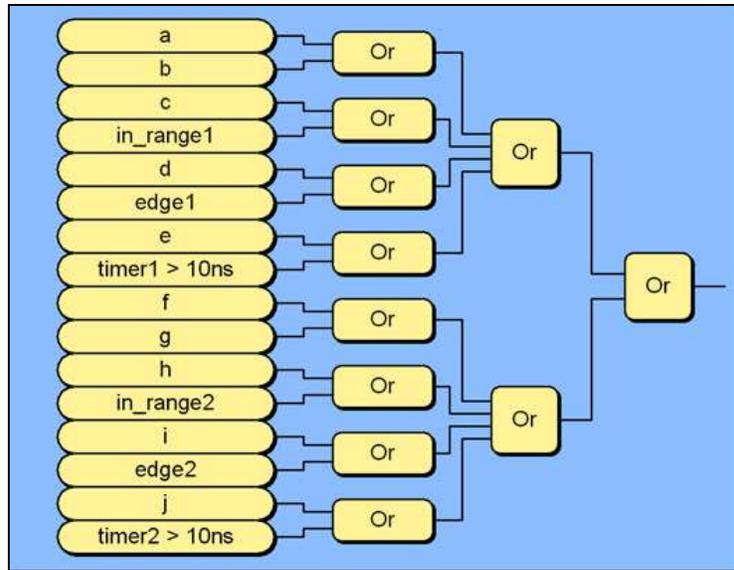


Figure 19 - Combination Trigger Sum

Each of the operations (the "OR" blocks shown above) are implemented using LUT memories. Therefore, they can be configured with any logical operation. Typical functions are: AND, NAND, OR, NOR, XOR, NXOR, A-only, B-only, ANY, or NOP.

In addition, the LUT memories can be used to invert the result of a trigger term, range, edge or timer detect (anything in the first column above). For example, allowing "a" or "not a", or "in_range" or "out of range" variations, etc...

There are three trigger sums per sequencer state, and sixteen trigger sequencer states, for a total of 48 trigger sums which need configuring.

7.1 Trigger Sum Inputs

As mentioned earlier, the inputs to the sum are trigger terms, range detectors, edge detectors and timers. Results of all inputs feed into the "pair" LUT's (which evaluate a pair of sum inputs each).

- Trigger Terms. 32-bit masked compares. Each term supplies a two bit result. LSB corresponds to indata[15:0]. MSB to indata[31:16]. Both bits must be asserted to indicate a hit.
- Range Detectors. The range detectors similarly supply two signals. LSB is the lower limit hit. MSB is upper limit hit. Both bits must be asserted to indicate a hit.
- Edge Detectors. Single hit/miss signal. Fed into both bits 2 & 3 of pair LUT.
- Timer. Single hit/miss signal. Fed into both bits 2 & 3 of pair LUT.

7.2 Trigger Sum Operations

Operation codes for the "pair" LUT's are a little confusing since terms & range detectors supply two signals each (both of which must be asserted for a hit).

| Sum Inputs | | | | AND | NAND | OR | NOR | XOR | NXOR | A | B |
|------------|---|---|---|-----|------|----|-----|-----|------|---|---|
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 1 |

Figure 20 - Pair LUT Operations

The mid LUT is straight forward. A pure 4-input logic gate:

| Sum Inputs | | | | AND | NAND | OR | NOR | XOR | NXOR |
|------------|---|---|---|-----|------|----|-----|-----|------|
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |

Figure 21 - Mid LUT Operations

The final LUT only has two inputs. It's configured as a 2-input logic gate:

| Sum Inputs | | | | AND | NAND | OR | NOR | XOR | NXOR |
|------------|---|---|---|-----|------|----|-----|-----|------|
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |

Figure 22 - Final LUT Operations

All LUT's can also be programmed to NOP (ie: always output zero), or ANY (always output one), or anything else you can think up.

7.3 Trigger Inverting

Inverting of trigger sum inputs is possible, but not obvious. Inverting of inputs must be incorporated into the LUT operation tables described in the previous section. Complicating things is how the "pair" LUT tables must combine two signals from the term & range detectors.

For example, see the OR/NOR function from the Pair LUT table below.

| Sum Inputs | | | | OR | NOR |
|------------|---|---|---|----|-----|
| B | A | | | | |
| 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 0 | 1 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 | 1 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 | 0 |

Inputs not inverted

| Sum Inputs | | | | OR | NOR |
|------------|---|---|---|----|-----|
| B | A | | | | |
| 0 | 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 0 | 1 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 | 0 |

Input "A" Inverted

| Sum Inputs | | | | OR | NOR |
|------------|---|---|---|----|-----|
| B | A | | | | |
| 0 | 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 0 | 1 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 1 | 1 | 0 |
| 1 | 1 | 0 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 |

Input "B" Inverted

The corresponding OR & NOR LUT table values are:

| | OR | NOR |
|--------------|--------|--------|
| Non Inverted | 0xF888 | 0x0777 |
| Invert "A" | 0xF777 | 0x0888 |
| Invert "B" | 0x8FFF | 0x7000 |
| Invert Both | 0x7FFF | 0x8000 |

To invert A, we simply exchange 0x7's with 0x8's (& visa versa). To invert B, we swap the most significant nibble with all lower nibbles.

7.4 Example (Trigger Sum Initialization)

```

#define OP_NOP 0
#define OP_ANY 1
#define OP_AND 2
#define OP_NAND 3
#define OP_OR 4
#define OP_NOR 5
#define OP_XOR 6
#define OP_NXOR 7
#define OP_A 8
#define OP_B 9

//
int pairvalue[] = {0x0000,0xFFFF,0x8000,0x7FFF,0xF888,0x0777,0x7888,0x8777,0x8888,0xF000};
int midvalue[] = {0x0000,0xFFFF,0x8000,0x7FFF,0xFFFE,0x0001,0x0116,0xFEE9,0xEEEE,0xFFF0};
int finalvalue[] = {0x0000,0xFFFF,0x0008,0x0007,0x000E,0x0001,0x0006,0x0009,0x0002,0x0004};

int get_pairvalue(int op, int invert_a, int invert_b)
{
    static int flipa[] = {0,0,0,0,0,0,0,8,7,0,0,0,0,0,0xF};
    int value = pairvalue[op];

    if (invert_a)
        value = (flipa[(value>>12)&0xF]<<12) | (flipa[(value>>8)&0xF]<<8) |
                (flipa[(value>>4)&0xF]<<4) | (flipa[value&0xF]);

    if (invert_b)
        value = ((value & 0xF)<<12) | ((value>>4) & 0xF0) |
                ((value>>8) & 0xF0) | ((value>>12) & 0xF);

    return (value);
}

void write_trigger_sum (
    int statenum, int stateterm,
    int invert_a, int invert_b, int invert_c, int invert_d, int invert_e,
    int invert_range1, int invert_edge1, int invert_timer1,
    int invert_f, int invert_g, int invert_h, int invert_i, int invert_j,
    int invert_range2, int invert_edge2, int invert_timer2,
    int op_ab, int op_c_range1, int op_d_edge1, int op_e_timer1,
    int op_fg, int op_h_range2, int op_i_edge2, int op_j_timer2,
    int op_mid1, int op_mid2, int op_final)
{
    int pv_ab = get_pairvalue(op_ab, invert_a, invert_b);
    int pv_c_range1 = get_pairvalue(op_c_range1, invert_c, invert_range1);
    int pv_d_edge1 = get_pairvalue(op_d_edge1, invert_d, invert_edge1);
    int pv_e_timer1 = get_pairvalue(op_e_timer1, invert_e, invert_timer1);
    int pv_fg = get_pairvalue(op_fg, invert_f, invert_g);
    int pv_h_range2 = get_pairvalue(op_h_range2, invert_h, invert_range2);
    int pv_i_edge2 = get_pairvalue(op_i_edge2, invert_i, invert_edge2);
    int pv_j_timer2 = get_pairvalue(op_j_timer2, invert_j, invert_timer2);

    write_select (0x40 + (statenum*4) + stateterm);
    write_chain (finalvalue[op_final]);
    write_chain ((midvalue[op_mid2]<<16) | midvalue[op_mid1]);
    write_chain ((pv_j_timer2<<16) | pv_i_edge2);
    write_chain ((pv_h_range2<<16) | pv_fg);
    write_chain ((pv_e_timer1<<16) | pv_d_edge1);
    write_chain ((pv_c_range1<<16) | pv_ab);
}

```

Figure 23 - Initialize Trigger Sum

8 Trigger Sequence States

Sequence states allow you to control how the analyzer evaluates and stores sampled data. Each state gives you control over data capturing, searching for a target (optionally multiple times), or branching to another sequence should a special condition occur. On a hit, the analyzer advances to the next sequence state. It can also trigger the analyzer controller & start/stop timers. With 16 sequence states, you can create complex protocol following triggers.

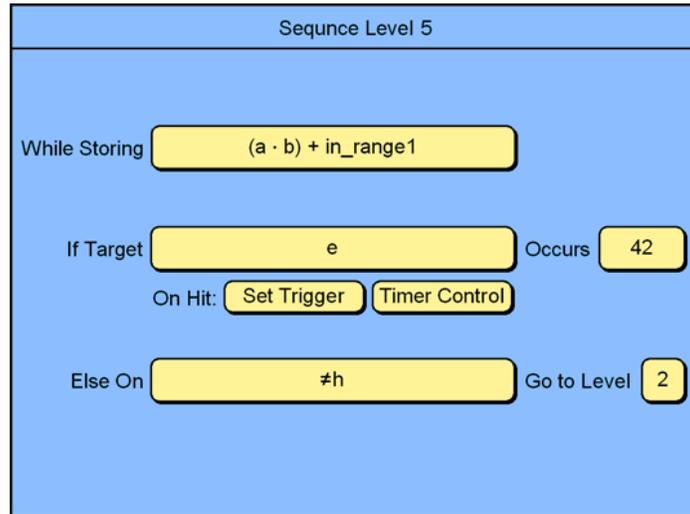


Figure 24 - Example Sequence Level

For each state you specify:

- Capture Term (sum of trigger terms, range checks, edge detect & timer hits)
- Hit Term (sum of trigger terms, range checks, edge detect & timer hits)
- Else Term (sum of trigger terms, range checks, edge detect & timer hits)
- Occurrence Count
- Else State
- Timer Control Flags (start/stop/clear for timers 1 & 2)
- Trigger Flag

The trigger sequencer itself is a relatively simple state machine. Pseudocode for it is:

```

if (sampled_input valid) then {
  if (capture_term valid) then
    capture data

  if (hit_term valid) then {
    increment hit_count
    if (hit_count == fsm_occurrence_count) then {
      update_timers
      reset hit_count
      if (fsm_trigger || last_state) then assert run to controller
      if (!last_state && !fsm_last_state) advance to next state
    }
  }
  otherwise if (else_term valid) then {
    reset hit_count
    go to state "fsm_else_level"
  }
}
    
```

Figure 25 - Trigger Sequencer State Machine

8.1 Example (Trigger Sequence Initialization)

```
#define TRIGSTATE_STATENUM_MASK 0xF
#define TRIGSTATE_OBTAIN_MASK 0x000FFFFF
#define TRIGSTATE_ELSE_BITOFS 20
#define TRIGSTATE_STOP_TIMER0 0x01000000
#define TRIGSTATE_STOP_TIMER1 0x02000000
#define TRIGSTATE_CLEAR_TIMER0 0x04000000
#define TRIGSTATE_CLEAR_TIMER1 0x08000000
#define TRIGSTATE_START_TIMER0 0x10000000
#define TRIGSTATE_START_TIMER1 0x20000000
#define TRIGSTATE_TRIGGER_FLAG 0x40000000
#define TRIGSTATE_LASTSTATE 0x80000000

void write_trigger_state (
    int statenum, // 0 to 15
    int last_state,
    int set_trigger,
    int start_timer, // bit0=timer1, bit1=timer2
    int stop_timer, // bit0=timer1, bit1=timer2
    int clear_timer, // bit0=timer1, bit1=timer2
    int else_state, // 0 to 15
    int obtain_count)
{
    write_select (statenum & TRIGSTATE_STATENUM_MASK);

    unsigned int value =
        ((else_state & TRIGSTATE_STATENUM_MASK)<<TRIGSTATE_ELSE_BITOFS) |
        (obtain_count & TRIGSTATE_OBTAIN_MASK);
    if (last_state) value |= TRIGGER_LASTSTATE;
    if (set_trigger) value |= TRIGSTATE_TRIGGER_FLAG;
    if (start_timer & 1) value |= TRIGSTATE_START_TIMER0;
    if (start_timer & 2) value |= TRIGSTATE_START_TIMER1;
    if (stop_timer & 1) value |= TRIGSTATE_STOP_TIMER0;
    if (stop_timer & 2) value |= TRIGSTATE_STOP_TIMER1;
    if (clear_timer & 1) value |= TRIGSTATE_CLEAR_TIMER0;
    if (clear_timer & 2) value |= TRIGSTATE_CLEAR_TIMER1;
    write_chain (value);
}
```

Figure 26 - Initialize Trigger Sequence

Appendix A - Advanced Trigger Config Registers

The client initializes advanced trigger LUT chains using the 0x9E *configuration-select* command, followed by one or more 0x9F *write-data* commands. Each data write adds 32-bits to the selected LUT chain.

Exceptions are the FSM state configs & timer limits, which store DWORD values directly. However, the client sees no difference (just more writes to trigger config regs).

Client Long Commands (command byte + four data bytes, LSB data byte first):

0x9E = ADVTRIG-Config-Select (LSB = chain#. other bytes reserved.)
0x9F = ADVTRIG-Write-Data

FSM State Data

| | | |
|------|----------------|-----------------|
| 0x00 | = FSM State 0 | (32-bit word's) |
| 0x01 | = FSM State 1 | |
| 0x02 | = FSM State 2 | |
| 0x03 | = FSM State 3 | |
| 0x04 | = FSM State 4 | |
| 0x05 | = FSM State 5 | |
| 0x06 | = FSM State 6 | |
| 0x07 | = FSM State 7 | |
| 0x08 | = FSM State 8 | |
| 0x09 | = FSM State 9 | |
| 0x0A | = FSM State 10 | |
| 0x0B | = FSM State 11 | |
| 0x0C | = FSM State 12 | |
| 0x0D | = FSM State 13 | |
| 0x0E | = FSM State 14 | |
| 0x0F | = FSM State 15 | |

LUTCHAIN's:

| | | |
|-----------|------------------|------------------|
| 0x20 | = trigger term a | (128 bit chains) |
| 0x21 | = trigger term b | |
| 0x22 | = trigger term c | |
| 0x23 | = trigger term d | |
| 0x24 | = trigger term e | |
| 0x25 | = trigger term f | |
| 0x26 | = trigger term g | |
| 0x27 | = trigger term h | |
| 0x28 | = trigger term i | |
| 0x29 | = trigger term j | |
| 0x30 | = range 1 lower | (512 bit chains) |
| 0x31 | = range 1 upper | (512 bit chains) |
| 0x32 | = range 2 lower | (512 bit chains) |
| 0x33 | = range 2 upper | (512 bit chains) |
| 0x34 | = edge 1 | (64 bit chains) |
| 0x35 | = edge 2 | (64 bit chains) |
| 0x38-0x39 | = timer 1 limit | (36-bit word's) |
| 0x3A-0x3B | = timer 2 limit | (36-bit word's) |

| | |
|-----------|------------------------|
| 0x40 | = state 0 hit-term |
| 0x41 | = state 0 else-term |
| 0x42 | = state 0 capture-term |
| 0x44-0x46 | = state 1 terms |
| 0x48-0x4A | = state 2 terms |
| 0x4C-0x4E | = state 3 terms |
| 0x50-0x52 | = state 4 terms |
| 0x54-0x56 | = state 5 terms |
| 0x58-0x5A | = state 6 terms |
| 0x5C-0x5E | = state 7 terms |
| 0x60-0x62 | = state 8 terms |
| 0x64-0x66 | = state 9 terms |
| 0x68-0x6A | = state 10 terms |
| 0x6C-0x6E | = state 11 terms |
| 0x70-0x72 | = state 12 terms |
| 0x74-0x76 | = state 13 terms |
| 0x78-0x7A | = state 14 terms |
| 0x7C-0x7E | = state 15 terms |

Appendix B - Example Code

```
#define write_select(value) write_long_command (0x9E, value)
#define write_chain(value) write_long_command (0x9F, value)

#define RANGE_XOR0 0xAAAA
#define RANGE_XOR1 0x5555
#define RANGE_NOP 0xFFFF

#define EDGE_RISE0 0x0A0A
#define EDGE_RISE1 0x00CC
#define EDGE_FALL0 0x5050
#define EDGE_FALL1 0x3300
#define EDGE_BOTH0 (EDGE_RISE0|EDGE_FALL0)
#define EDGE_BOTH1 (EDGE_RISE1|EDGE_FALL1)
#define EDGE_NEITHER0 (~EDGE_BOTH0 & 0xFFFF)
#define EDGE_NEITHER1 (~EDGE_BOTH1 & 0xFFFF)

#define TRIGSTATE_STATENUM_MASK 0xF
#define TRIGSTATE_OBTAIN_MASK 0x000FFFFF
#define TRIGSTATE_ELSE_BITOFS 20
#define TRIGSTATE_STOP_TIMER0 0x01000000
#define TRIGSTATE_STOP_TIMER1 0x02000000
#define TRIGSTATE_CLEAR_TIMER0 0x04000000
#define TRIGSTATE_CLEAR_TIMER1 0x08000000
#define TRIGSTATE_START_TIMER0 0x10000000
#define TRIGSTATE_START_TIMER1 0x20000000
#define TRIGSTATE_TRIGGER_FLAG 0x40000000
#define TRIGSTATE_LASTSTATE 0x80000000

#define OP_NOP 0
#define OP_ANY 1
#define OP_AND 2
#define OP_NAND 3
#define OP_OR 4
#define OP_NOR 5
#define OP_XOR 6
#define OP_NXOR 7
#define OP_A 8
#define OP_B 9

//
// Setup LUT's for Trigger Term
//
// Inputs are 32-bit target & mask for comparing against captured analyzer data.
// If a mask bit is set, the corresponding target bit participates in the trigger.
//
void write_term (int term_number, unsigned int target, unsigned int mask)
{
    unsigned int bitmask=1;
    unsigned int lutvalue0=0;
    unsigned int lutvalue1=0;
    unsigned int lutvalue2=0;
    unsigned int lutvalue3=0;

    for (int i=0; i<16; i=i+1) {
        if (((i ^ (target & 0xF)) & (mask & 0xF))==0) lutvalue0 |= bitmask;
        if (((i ^ ((target>>4) & 0xF)) & ((mask>>4) & 0xF))==0) lutvalue0 |= (bitmask<<16);
        if (((i ^ ((target>>8) & 0xF)) & ((mask>>8) & 0xF))==0) lutvalue1 |= bitmask;
        if (((i ^ ((target>>12) & 0xF)) & ((mask>>12) & 0xF))==0) lutvalue1 |= (bitmask<<16);
        if (((i ^ ((target>>16) & 0xF)) & ((mask>>16) & 0xF))==0) lutvalue2 |= bitmask;
        if (((i ^ ((target>>20) & 0xF)) & ((mask>>20) & 0xF))==0) lutvalue2 |= (bitmask<<16);
        if (((i ^ ((target>>24) & 0xF)) & ((mask>>24) & 0xF))==0) lutvalue3 |= bitmask;
        if (((i ^ ((target>>28) & 0xF)) & ((mask>>28) & 0xF))==0) lutvalue3 |= (bitmask<<16);
        bitmask <<= 1;
    }

    // write data into LUT serial chain. MSB must goes in first. Total of 128 bits.
    write_select (0x20 + (term_number%10));
    write_chain (lutvalue3);
    write_chain (lutvalue2);
    write_chain (lutvalue1);
    write_chain (lutvalue0);
}
```

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```
//
// Setup LUT's for Range Detectors
//
// Inputs:
//   rangesel: 0=range1-lower, 1=range1-upper, 2=range2-lower, 3=range2-upper
//   target:   Limit value.
//   mask:     Indicate which bits should participate in range compare.
//
void write_range (int rangesel, unsigned int target, unsigned int mask)
{
    unsigned int value;
    unsigned int lutvalue=0;
    int i;

    write_select(0x30+(rangesel&3));

    // Count # of bits in mask...
    int bitcount=0;
    unsigned int bitmask=1;
    for (i=0; i<32; i++, bitmask<<=1)
        if (mask & bitmask)
            bitcount++;

    // Prepare target value...
    if (rangesel & 1)
        value = ~target; // upper range target
    else value = ~(target-1); // lower range value

    // Push MSB of target into bit 31...
    value <<= (32-bitcount);

    // Generate & program LUT values. Total of 512 bits.
    for (i=0; i<16; i=i+1) {
        if (((mask>>31)&1)==0)
            lutvalue = RANGE_NOP;
        else {
            lutvalue = ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
            value <<= 1;
        }
        mask <<= 1;
        lutvalue <<= 16;

        if (((mask>>31)&1)==0)
            lutvalue |= RANGE_NOP;
        else {
            lutvalue |= ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
            value <<= 1;
        }
        mask <<= 1;

        write_chain(lutvalue);
    }
}

//
// Setup LUT's for Edge Detectors.
// Each LUT handles two bits of input. Bits 0 & 1 are input. Bits 2 & 3 delayed input.
//
void write_edge (int edgesel,
                unsigned int rising_edge, unsigned int falling_edge, unsigned int neither_edge)
{
    write_select (0x34 + (edgesel&1));

    unsigned int lutvalue=0;
    unsigned int bitmask = 0x80000000;
    for (int i=0; i<16; i=i+1) {
        // Evaluate indata bit1...
        if (neither_edge && bitmask)
            lutvalue |= EDGE_NEITHER1;
        else {
            if (rising_edge & bitmask) lutvalue |= EDGE_RISE1;
            if (falling_edge & bitmask) lutvalue |= EDGE_FALL1;
        }
        bitmask >>= 1;

        // Evaluate indata bit0...
        if (neither_edge && bitmask)
            lutvalue |= EDGE_NEITHER0;
        else {
            if (rising_edge & bitmask) lutvalue |= EDGE_RISE0;
        }
    }
}
```

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```
    if (falling_edge & bitmask) lutvalue |= EDGE_FALL0;
}
bitmask >>= 1;

if ((i&1)==0)
    lutvalue <<= 16;
else {
    write_chain (lutvalue); // write total of 256 bits
    lutvalue = 0;
}
}
}

//
// Setup trigger timers...
//
void write_trigger_limit (int timer_sel, uint64_t value)
{
    write_select (0x38 + (timer_sel&1)*2);
    write_chain (value & 0xFFFFFFFF);
    write_select (0x39 + (timer_sel&1)*2);
    write_chain (value>>32);
}

//
// Setup trigger state term combinational sum. Operations for
// all 11 fields are merged & written. Uses table based lookups.
//
//
//
int pairvalue[] = {NOP ANY AND NAND OR NOR XOR NXOR A B};
int midvalue[] = {0x0000,0xFFFF,0x8000,0x7FFF,0xF888,0x0777,0x7888,0x8777,0x8888,0xF000};
int finalvalue[] = {0x0000,0xFFFF,0x0008,0x0007,0x000E,0x0001,0x0006,0x0009,0x0002,0x0004};

int get_pairvalue(int op, int invert_a, int invert_b)
{
    static int flipa[] = {0,0,0,0,0,0,0,8,7,0,0,0,0,0,0,0,0xF};
    int value = pairvalue[op];

    if (invert_a)
        value = (flipa[(value>>12)&0xF]<<12) | (flipa[(value>>8)&0xF]<<8) |
            (flipa[(value>>4)&0xF]<<4) | (flipa[value&0xF]);

    if (invert_b)
        value = ((value & 0xF)<<12) | ((value>>4) & 0xF00) |
            ((value>>8) & 0xF0) | ((value>>12) & 0xF);

    return (value);
}

void write_trigger_sum (
    int statenum, int stateterm,
    int invert_a, int invert_b, int invert_c, int invert_d, int invert_e,
    int invert_range1, int invert_edge1, int invert_timer1,
    int invert_f, int invert_g, int invert_h, int invert_i, int invert_j,
    int invert_range2, int invert_edge2, int invert_timer2,
    int op_ab, int op_c_range1, int op_d_edge1, int op_e_timer1,
    int op_fg, int op_h_range2, int op_i_edge2, int op_j_timer2,
    int op_mid1, int op_mid2, int op_final)
{
    int pv_ab = get_pairvalue(op_ab, invert_a, invert_b);
    int pv_c_range1 = get_pairvalue(op_c_range1, invert_c, invert_range1);
    int pv_d_edge1 = get_pairvalue(op_d_edge1, invert_d, invert_edge1);
    int pv_e_timer1 = get_pairvalue(op_e_timer1, invert_e, invert_timer1);
    int pv_fg = get_pairvalue(op_fg, invert_f, invert_g);
    int pv_h_range2 = get_pairvalue(op_h_range2, invert_h, invert_range2);
    int pv_i_edge2 = get_pairvalue(op_i_edge2, invert_i, invert_edge2);
    int pv_j_timer2 = get_pairvalue(op_j_timer2, invert_j, invert_timer2);

    write_select (0x40 + (statenum*4) + stateterm);
    write_chain (finalvalue[op_final]);
    write_chain ((midvalue[op_mid2]<<16) | midvalue[op_mid1]);
    write_chain ((pv_j_timer2<<16) | pv_i_edge2);
    write_chain ((pv_h_range2<<16) | pv_fg);
    write_chain ((pv_e_timer1<<16) | pv_d_edge1);
    write_chain ((pv_c_range1<<16) | pv_ab);
}
}
```

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```
//  
// Setup trigger state. Specify trigger mode, timer controls, "else" state number,  
// and necessary "obtain" count (number of hits needed before full hit observed).  
//  
void write_trigger_state (  
    int statenum, // 0 to 15  
    int last_state,  
    int set_trigger,  
    int start_timer, // bit0=timer1, bit1=timer2  
    int stop_timer, // bit0=timer1, bit1=timer2  
    int clear_timer, // bit0=timer1, bit1=timer2  
    int else_state, // 0 to 15  
    int obtain_count)  
{  
    write_select (statenum & TRIGSTATE_STATENUM_MASK);  
  
    unsigned int value =  
        ((else_state & TRIGSTATE_STATENUM_MASK)<<TRIGSTATE_ELSE_BITOFS) |  
        (obtain_count & TRIGSTATE_OBTAIN_MASK);  
    if (last_state) value |= TRIGSTATE_LASTSTATE;  
    if (set_trigger) value |= TRIGSTATE_TRIGGER_FLAG;  
    if (start_timer & 1) value |= TRIGSTATE_START_TIMER0;  
    if (start_timer & 2) value |= TRIGSTATE_START_TIMER1;  
    if (stop_timer & 1) value |= TRIGSTATE_STOP_TIMER0;  
    if (stop_timer & 2) value |= TRIGSTATE_STOP_TIMER1;  
    if (clear_timer & 1) value |= TRIGSTATE_CLEAR_TIMER0;  
    if (clear_timer & 2) value |= TRIGSTATE_CLEAR_TIMER1;  
    write_chain (value);  
}
```

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